

# On the strength of Hindman's Theorem for bounded sums of unions

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# Outline

- 1 Hindman's Finite Sums Theorem
- 2 Bounded Sums
- 3 Weak Yet Strong Principles
- 4 From Hindman to Ramsey
- 5 Other variants

# Hindman's Finite Sums Theorem

## Theorem (Hindman, 1972)

*Whenever the positive integers are colored in finitely many colors there is an infinite set such that all non-empty finite sums of distinct elements drawn from that set have the same color.*

- Original proof is combinatorial but intricate.
- Later proofs are simpler but use strong methods (ultrafilters or ergodic theory).

## Question, '80s

What is the strength of Hindman's Theorem?

# Measures of Strength

$$\text{HT}_k = \forall \underbrace{c : \mathbf{N} \rightarrow k}_{\text{instance}} \exists \underbrace{X \subseteq \mathbf{N}}_{\text{solution}} (|X| = \aleph_0 \text{ and } FS(X) \text{ is mono})$$

$$\text{HT} = \forall k \text{HT}_k$$

- Reverse Mathematics: **provability** in the systems

$$\text{RCA}_0, \text{WKL}_0, \text{ACA}_0, \text{ACA}'_0, \text{ACA}^+_0, \dots$$

or (mutual) implications over the base theory  $\text{RCA}_0$ .

- Computable Mathematics: **complexity of solutions** for **computable instances**.
- RM and CM: computable **reducibility** to/from other principles.

# Lower Bound on Hindman's Theorem

$$\text{HT} \geq \emptyset^{(1)}, \text{RT}_2^3, \text{ACA}_0$$

Theorem (Blass, Hirst, Simpson 1987)

- 1 *Some computable (resp. computable in  $X$ ) 2-coloring of  $\mathbf{N}$  admits only solutions to  $\text{HT}_2$  that compute  $\emptyset^{(1)}$  (resp.  $X'$  – the jump of  $X$ ).*
- 2  $\text{RCA}_0 + \text{HT}_2 \vdash \text{ACA}_0$ .

- Proof is by coding of the Halting Set and formalizes in  $\text{RCA}_0$ .
- Uses the notion of gap, the interval between two successive exponents of a number in base 2.

# Upper Bound on Hindman's Theorem

$$\text{ACA}_0^+, \emptyset^{(\omega+1)} \geq \text{HT}$$

Theorem (Blass, Hirst, Simpson 1987)

- 1 Any finite computable (resp. computable in  $X$ ) coloring of  $\mathbf{N}$  admits a solution to HT computable in  $\emptyset^{(\omega+1)}$  (resp. in  $X^{(\omega+1)}$ ).
- 2  $\text{ACA}_0^+ \vdash \text{HT}$ .

- $\text{ACA}_0^+$  is  $\text{ACA}_0$  plus  $\forall X \exists Y (Y = X^{(\omega)})$ .
- Proof is by analyzing the original proof by Hindman.
- Ultrafilter and ergodic proofs give worse bounds (so far).

# Bounded Sums

## Question (Blass, 2005)

Does the complexity of HT grow with the length of the sums?

- Is it the case that longer sums require more jumps?
- $FS(X)$  = sums of finitely many distinct elements of  $X$ .
- $FS^{\leq n}(X)$  = sums of  $1, 2, \dots, n$  distinct elements of  $X$ .
- $HT_k^{\leq n}$  = the restriction of HT to  $k$  colors and sums of length  $\leq n$ .

$$HT_k^{\leq n}, HT^{\leq n}$$

# Lower Bounds for bounded sums

$$\text{HT}^{\leq 3} \geq \emptyset^{(1)}, \text{RT}_2^3, \text{ACA}_0$$

Theorem (Dzhafarov, Jockusch, Solomon, Westrick, 2017)

- 1  $\text{RCA}_0 + \text{HT}_3^{\leq 3} \vdash \text{ACA}_0$ .
- 2  $\text{RCA}_0 \not\vdash \text{HT}_2^{\leq 2}$ , and  $\text{RCA}_0 + \text{RT}^1 + \text{HT}_2^{\leq 2} \vdash \text{SRT}_2^2$ .

- $\text{SRT}_2^2$  is the Stable Ramsey's Theorem ( $\text{WKL}_0 \not\vdash \text{SRT}_2^2$ ).
- Proof of (1): modification of Blass-Hirst-Simpson's argument.
- Proof of (2): Given a  $\Delta_2^0$ -set  $A$  define a coloring all of whose solutions compute an infinite subset of  $A$  or an infinite set disjoint from  $A$ . Formalization requires  $\text{RT}^1$  (eq.  $B\Sigma_2^0$ ).



# Upper bounds for bounded sums?

Question (Hindman, Leader and Strauss, 2003)

Is there a proof that whenever  $\mathbf{N}$  is finitely coloured there is a sequence  $x_1, x_2, \dots$  such that all  $x_i$  and all  $x_i + x_j$  ( $i \neq j$ ) have the same colour, that does not also prove the Finite Sums Theorem?

- Does  $\text{HT}^{\leq 2}$  imply HT over  $\text{RCA}_0$ ?
- Can we upper bound  $\text{HT}^{\leq 2}$  below  $\text{ACA}_0^+$ ?
- Are there natural Hindman-type principles with:
  - 1 Non-trivial lower bounds, and
  - 2 Upper bounds strictly below HT?
- We call such principles Weak Yet Strong.

# A brute force proof using Ramsey

Given  $c : \mathbf{N} \rightarrow 2$ ,

1. Use  $RT_2^1$  on  $\mathbf{N}$  wrt  $c$  to get an infinite homset  $H_1$ .
2. Use  $RT_2^2$  on  $H_1$  wrt  $f_2(x, y) := c(x + y)$  to fix the color of sums of length 2 on an infinite  $H_2 \subseteq H_1$ .
- ...
- k. Use  $RT_2^k$  on  $H_{k-1}$  wrt  $f_k(x_1, \dots, x_k) := c(x_1 + \dots + x_k)$  to fix the color of sums of length  $k$  on an infinite  $H_k \subseteq H_{k-1}$ .

This induces a coloring  $d : [1, k] \rightarrow 2$ , where  $d(i)$  is the  $c$ -color of sums of length  $i$  from  $H_k$ .

If  $k$  is large, then  $d$  has some interesting homogeneous set!

E.g. if  $k \geq 6$  then by Schur's Theorem there exists  $a, b > 0$  such that

$$d(a) = d(b) = d(a + b).$$

# Hindman-Schur Theorem

- $FS^A(X)$ : sums of  $j$ -many distinct elements of  $X$  for any  $j \in A$ .
- **Hindman-Schur Theorem**: Whenever the positive integers are colored in two colors **there exist** positive integers  $a, b$  and an infinite set  $H$  such that  $FS^{\{a,b,a+b\}}(H)$  is monochromatic.

## Theorem (C., 2017)

*Hindman-Schur Theorem is provable in  $ACA_0$ .*

- A host of similar Hindman-type theorems based on different finite combinatorial principles (e.g., Van Der Waerden, Folkman, etc.).
- All provable in  $ACA_0$ .
- What about lower bounds?

# Hindman-Schur with apartness

The Blass-Hirst-Simpson's lower bound proof works, **if we impose that the solution set satisfies the following Apartness Condition, for  $t = 2$ .**

## Definition ( $t$ -Apartness)

Fix a base  $t \geq 2$ . A set  $X \subseteq \mathbf{N}$  satisfies the  $t$ -apartness condition if

$$x < x' \Rightarrow \mu_t(x) < \lambda_t(x').$$

$\lambda_t(x)$  = least exponent in base  $t$  representation of  $n$ .

$\mu_t(x)$  = maximal exponent in base  $t$  representation of  $n$ .

- $P$  with  $t$ -apartness =  $P$  with  $t$ -apartness on the solution set.

## Theorem (C., Kołodziejczyk, Lepore, Zdanowski, 2017)

*Hindman-Schur with 2-apartness is equivalent to  $ACA_0$  (over  $RCA_0$ ).*

# The Apartness Condition

Imposing apartness is a **self-strengthening** of Hindman's Theorem:

$$\text{RCA}_0 \vdash \text{HT} \equiv \text{HT with apartness.}$$

For restricted versions we have the following:

Proposition (C., Kołodziejczyk, Lepore, Zdanowski, 2017)

$\text{RCA}_0 + \text{HT}_{2k}^{\leq n} \vdash \text{HT}_k^{\leq n}$  with 3-apartness.

**Proof:** Give  $c : \mathbf{N} \rightarrow 2$ , let  $d : \mathbf{N} \rightarrow 4$ :

$$d(n) := \begin{cases} c(n) & \text{if } n = 3^t + \dots, \\ 2 + c(n) & \text{if } n = 2 \cdot 3^t + \dots \end{cases}$$

If  $FS^{\leq 2}(H)$  is monochromatic for  $d$  then:

- 1 all elements have same first coefficient. Then:
- 2 no two elements of  $H$  can have the same first exponent.

# Bounded Unions

Let  $\text{FUT}_k^{\leq n}$  and  $\text{FUT}_k^{=n}$  the versions of Hindman's Theorem in terms of unions instead of sums.

## Proposition

*For each  $n, kt \geq 2$ ,  $\text{HT}_k^{\leq n}$  with  $t$ -apartness is equivalent to  $\text{FUT}_k^{\leq n}$  over  $\text{RCA}_0$ . Moreover, these principles are mutually strongly computably reducible. The same equivalences hold for  $\text{HT}_k^{=n}$  with  $t$ -apartness and  $\text{FUT}_k^{=n}$ .*

# Restricted Hindman and Polarized Ramsey

Recall that Dzhafarov et alii proved

$$\text{RCA}_0 + \text{HT}^{\leq 2} + \text{RT}^1 \vdash \text{SRT}_2^2$$

We improve by showing that

$$\text{RCA}_0 + \text{HT}^{\leq 2} \vdash \text{IPT}_2^2$$

Definition (Dzhafarov and Hirst, 2011)

$\text{IPT}_2^2$ : For all  $f : [\mathbf{N}]^2 \rightarrow 2$  there exists a pair of infinite sets  $(H_1, H_2)$  such that all increasing pairs  $\{x_1, x_2\}$  with  $x_i \in H_i$  get the same  $f$ -color.

$$\text{RT}_2^2 \geq \text{IPT}_2^2 > \text{SRT}_2^2$$

## Restricted Hindman and Polarized Ramsey

In fact we get that  $\text{IPT}_2^2$  is **strongly computably reducible** to  $\text{HT}_4^{\leq 2}$ : any  $f : [\mathbf{N}]^2 \rightarrow 2$  of  $\text{IPT}_2^2$  computes an instance  $c : \mathbf{N} \rightarrow 2$  of  $\text{HT}_4^{\leq 2}$  s.t. any solution to  $\text{HT}_4^{\leq 2}$  for  $c$  computes a solution to  $\text{IPT}_2^2$  for  $f$ .

### Theorem (C., 2017)

$\text{RCA}_0 + \text{HT}_4^{\leq 2} \vdash \text{IPT}_2^2$ . Moreover,  $\text{IPT}_2^2 \leq_{\text{sc}} \text{HT}_4^{\leq 2}$ .

- $\text{HT}_k^{\overline{n}}$  = restriction of  $\text{HT}_k$  to sums of *exactly*  $n$  elements.

In fact we show:

### Theorem (C., 2017)

$\text{RCA}_0 + \text{HT}_2^{\overline{2}}$  with  $t$ -apartness  $\vdash \text{IPT}_2^2$ . Moreover,  $\text{IPT}_2^2 \leq_{\text{sc}} \text{HT}_2^{\overline{2}}$  with  $t$ -apartness.

- N.B.  $\text{RT}_2^2$  proves  $\text{HT}_2^{\overline{2}}$  with  $t$ -apartness.



# $\text{IPT}_2^2 \leq_{\text{sc}} \text{HT}_2^2$ with 2-apartness

Given  $f : [\mathbf{N}]^2 \rightarrow 2$ , let  $g : \mathbf{N} \rightarrow 2$ :

$$g(n) := \begin{cases} 0 & \text{if } n = 2^t, \\ f(\lambda(n), \mu(n)) & \text{otherwise.} \end{cases}$$

Let  $H = \{h_1 < h_2 < h_3 < \dots\}$  be an infinite and 2-apart set such that  $g$  is constant on  $FS^{=2}(H)$ . Then

$$\lambda(h_1) \leq \mu(h_1) < \lambda(h_2) \leq \mu(h_2) < \lambda(h_3) \leq \mu(h_3) < \dots$$

So if

$$H_1 := \{\lambda(h_1), \lambda(h_3), \lambda(h_5), \dots, \}$$

$$H_2 := \{\mu(h_2), \mu(h_4), \mu(h_6), \dots, \}$$

Then  $(H_1, H_2)$  is a solution to  $\text{IPT}_2^2$  for  $f$ .

# Sums of length 2 and $ACA_0$

$$HT^{\leq 2} \geq \emptyset^{(1)}, RT_2^3, ACA_0$$

Recall that Dzhafarov et alii proved

$$RCA_0 + HT^{\leq 3} \vdash ACA_0.$$

Theorem (C., Kołodziejczyk, Lepore, Zdanowski, 2017)

$$RCA_0 + HT^{\leq 2} \vdash ACA_0.$$

Proposition (C., Kołodziejczyk, Lepore, Zdanowski, 2017)

For  $t \geq 2$ ,  $RCA_0 + HT_2^{\leq 2}$  with  $t$ -apartness  $\vdash ACA_0$ .

## $HT_2^{\leq 2}$ with apartness implies $ACA_0$

Let  $f : \mathbf{N} \rightarrow \mathbf{N}$  be 1:1. Let  $n = 2^{n_0} + \dots + 2^{n_r}$ , ( $n_0 < \dots < n_r$ ). Consider

$$f \upharpoonright [0, n_0), f \upharpoonright [n_0, n_1), \dots, f \upharpoonright [n_{r-1}, n_r).$$

Call  $j \leq r$  **important in**  $n$  iff some value of  $f \upharpoonright [n_{j-1}, n_j)$  is below  $n_0$ . ( $n_{-1} := 0$ ).

$$c(n) := \text{parity of the number of important } js \text{ in } n.$$

Let  $H$  be infinite, 2-apart and  $FS^{\leq 2}(H)$  mono.

**Claim:** for each  $n \in H$  and each  $x < \lambda(n)$ ,

$$x \in \text{rg}(f) \text{ if and only if } x \in \text{rg}(f \upharpoonright \mu(n)).$$

Gives a computable definition of  $\text{rg}(f)$ : given  $x$ , find the smallest  $n \in H$  such that  $x < \lambda(n)$  and check whether  $x$  is in  $\text{rg}(f \upharpoonright \mu(n))$ .

# $HT_k^{\overline{n}}$ with apartness and $ACA_0$

By improving the proof we get the following:

Proposition (C., Kołodziejczyk, Lepore, Zdanowski, 2017)

For every  $t \geq 2$ ,  $RCA_0 + HT_2^{\overline{3}}$  with  $t$ -apartness  $\vdash ACA_0$ .

- Therefore  $\{HT_k^{\overline{n}}$  with 2-apartness ;  $n \geq 3, k \geq 2\}$  is a weak yet strong family.

Corollary (C., Kołodziejczyk, Lepore, Zdanowski, 2017)

For every  $n \geq 3$  and  $k \geq 2$ ,

$$HT_k^{\overline{n}} \text{ with 2-apartness} \equiv ACA_0$$

over  $RCA_0$ .

# Open Problems

- Can we upper bound  $HT_2^{\leq 2}$  strictly below  $ACA_0^+$ ?
- Is  $HT_2^{\leq 2}$  provable in  $ACA_0$ ?
- Do colors matter? How?
- Does apartness increase strength in the bounded cases?
- Which implications are witnessed by reductions? E.g. Does  $IPT_2^3 \leq_{sc} HT_2^{\leq 3}$ ?

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