

Caristi's fixed point theorem and strong systems of arithmetic

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Fixed point theorems in analysis

Theorem (Banach, 1922)

Let \mathcal{X} be a complete metric space and $f: \mathcal{X} \rightarrow \mathcal{X}$ be a *contraction*; that is, there is $\rho < 1$ such that $d(f(x), f(y)) < \rho \cdot d(x, y)$ for all $x, y \in X$. Then, there is $x_* \in X$ such that $f(x_*) = x_*$.

Fixed point theorems in analysis

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Theorem (Brouwer, 1910)

Let D be a disk in \mathbb{R}^n and $f: D \rightarrow D$ be continuous. Then, there is $x_* \in X$ such that $f(x_*) = x_*$.

Caristi systems and Caristi's fixed point theorem

Definition

A **Caristi system** is a triple (\mathcal{X}, V, f) , where

- \mathcal{X} is a complete separable metric space,
- $V: \mathcal{X} \rightarrow \mathbb{R}^{\geq 0}$ is a **lower semi-continuous** function, and
- $f: \mathcal{X} \rightarrow \mathcal{X}$ is an **arbitrary function**,

such that

$$\forall x \in \mathcal{X} \left(d(x, f(x)) \leq V(x) - V(f(x)) \right).$$

Theorem (Caristi, 1976)

If (\mathcal{X}, V, f) is a Caristi system, then f has a fixed point.

- Henceforth, a **metric space** is a complete separable metric space.
- We call the V in a Caristi system a **potential**.

Proofs of Caristi's theorem

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- 2 Proof by EVP.

Theorem (Ekeland, 1974)

Every lower semi-continuous function $V: \mathcal{X} \rightarrow \mathbb{R}_{\geq 0}$ has a *critical point*, i.e. a point $x_* \in \mathcal{X}$ such that

$$\forall y \in \mathcal{X} \left(d(x_*, y) \leq V(x_*) - V(y) \rightarrow y = x_* \right)$$

The Big Five subsystems of second-order arithmetic

- **Two sorts:** natural numbers, sets of naturals.
- Real numbers, infinite trees, etc. can all be coded in SOA.
- All systems have Σ_1^0 -induction, elementary arithmetical axioms.

RCA₀ Computable sets exist (Δ_1^0 comprehension)

WKL₀ RCA₀ + “every infinite binary tree has an infinite path”

ACA₀ $\{n \in \mathbb{N} : \varphi(n)\}$ exists, φ arithmetical

ATR₀ Transfinite recursion for arithmetical formulas.

Π_1^1 -CA₀ $\{n \in \mathbb{N} : \forall X \subseteq \mathbb{N} \varphi(n, X)\}$ exists, φ arithmetical

The strength of continuous Caristi's theorem

Theorem (F-D S T Y)

The following are equivalent over RCA_0 .

- 1 ACA_0 .
- 2 *Caristi's theorem for continuous functions.*
- 3 *Caristi's theorem for continuous potentials and continuous functions.*

Proof. (1 \rightarrow 2). By EVP.

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(3 \rightarrow 1). Use the fact that ACA_0 is equivalent to the statement that every decreasing sequence of positive reals has an infimum.

Compactness and Caristi's theorem

Theorem (F-D S T Y)

For **compact** metric spaces \mathcal{X} :

- Caristi for l.s.c. V and continuous f is equivalent to WKL_0 .
- Caristi for continuous V and continuous f is equivalent to WKL_0 .

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Baire functions:

- Continuous functions are Baire class 0
- ω -limits of Baire class $< \xi$ functions are Baire class ξ .

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Theorem (F-D S T Y)

For **compact** metric spaces \mathcal{X} :

- Caristi for l.s.c. V and Baire class 1 f is equivalent to ACA_0 .

Caristi vs. ATR_0

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Caristi for Baire class 1 functions $f: \mathcal{X} \rightarrow \mathcal{X}$ (with arbitrary \mathcal{X} and l.s.c. V) implies ATR_0 .

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Facts:

- ① $\varphi\text{-TR} := \forall \prec \left(\text{WO}(\prec) \rightarrow \exists Z \forall \xi \forall n (n \in Z_\xi \leftrightarrow \varphi(n, Z_{\prec\xi})) \right)$
- $Z_\xi = \{m \in \mathbb{N} : \langle m, \xi \rangle \in Z\}$,
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- $\text{ATR}_0 \equiv \text{RCA}_0 + \Sigma_1^0\text{-TR}$.

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 - $Z_{\prec\xi} = \{m \in \mathbb{N} : \exists \zeta \prec \xi \langle m, \zeta \rangle \in Z\}$.
- ② $ATR_0 \equiv RCA_0 + \Sigma_1^0\text{-}TR$.
- ③ If φ is Σ_1^0 , there is a sequence of trees $(T_\xi)_{\xi \in \mathbb{N}}$ such that any path g through T_ξ codes Z_ξ .

Caristi for Baire class 1 F implies ATR_0 .

Proof idea. Fix $(T_i)_{i < \omega}$. We define a Caristi system (\mathcal{X}, F, V) as follows.

- \mathcal{X} be the set of sequences of paths $(g_i)_{i < \omega}$.
- $F(\vec{g}) = \lim_{n \rightarrow \omega} F_n$, where $F_n(\vec{g})$ replaces g_ξ by an ' n -approximation' of the path through T_ξ , whenever:
 - ① $g_\xi(n) \notin T_\xi$, and
 - ② $\xi < n$ is the \prec -minimum satisfying 1.
- $V(\vec{g}) = \sum \{2^{-n} : g_n \text{ is not a path through } T_i\}$.

This defines a Caristi system, whose fixed point $(g_i^*)_{i < \omega}$ such that each g_i^* is a path through T_i . □

Leftmost paths

Theorem (Marcone)

$\Pi_1^1\text{-CA}_0$ is equivalent to the statement “every ill-founded tree $T \subseteq \mathbb{N}^{<\mathbb{N}}$ has a leftmost path.”

Definition (Towsner)

- The **transfinite leftmost path principle** states that if $T \subseteq \mathbb{N}^{<\mathbb{N}}$ is ill-founded and α is a well-order, then there is a path f^* through T such that no path through T is both $\Sigma_\alpha^{T \oplus f^*}$ and to the left of f^* .
- TLPP_0 is RCA_0 plus the transfinite leftmost path principle.

TLPP_0 is strictly between ATR_0 and $\Pi_1^1\text{-CA}_0$.

Caristi's fixed point theorem for Baire functions

Theorem (F-D S T Y)

Caristi for Baire functions $f: \mathcal{X} \rightarrow \mathcal{X}$ (with arbitrary \mathcal{X} and l.s.c. V) is equivalent to TLPP_0 .

Thus in the **general case**:

- Caristi is equivalent to TLPP_0
- Ekeland is equivalent to $\Pi_1^1\text{-CA}_0$

Strength of Caristi's proof

Recall: Caristi's proof relies on uncountable Caristi sequences.

Definition

Fix a Caristi system (\mathcal{X}, V, f) .

A **Caristi sequence (from x_0)** is a well-order (L, \prec) and a sequence $(x_\ell : \ell \in L) \subseteq \mathcal{X}$ such that

$$\left\{ \begin{array}{l} x_{\min L} = x_0 \\ x_{S(\ell)} = f(x_\ell) \\ x_\ell = \lim_{k < \ell} x_k \end{array} \right. \quad (\ell \in \text{Lim})$$

Maximal sequences

Call a Caristi sequence **proper** if the x_ℓ 's are all distinct.

Lemma (Maximal sequence principle)

Given a Caristi system (\mathcal{X}, V, f) with f arithmetical and $x_0 \in \mathcal{X}$, there is a proper Caristi sequence with no strict, proper extensions.

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Proof of Caristi's theorem by MSP: Any maximal sequence must have a last element, which is a fixed point of f . □

The closed orbit principle

Closed orbit: If (\mathcal{X}, f) is a dynamical system and $x \in \mathcal{X}$, the **closed orbit** of x is the least topologically closed, f -closed set \mathcal{O}^* such that $x \in \mathcal{O}^*$.

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Theorem (Banach fixed point theorem 2.0)

If \mathcal{X} is a metric space and $f: \mathcal{X} \rightarrow \mathcal{X}$ is a contraction, then for any $x \in \mathcal{X}$, the closed orbit of \mathcal{X} has a unique fixed point.

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Lemma (Closed orbit principle)

For every Caristi system (\mathcal{X}, V, f) with f arithmetical and every $x_0 \in \mathcal{X}$, there is a \subseteq -least closed set \mathcal{O}^ such that $x_0 \in \mathcal{O}^*$ and $(\forall x \in \mathcal{X})(x \in \mathcal{O}^* \rightarrow f(x) \in \mathcal{O}^*)$.*

Inflationary fixed points

Definition

The **arithmetical inflationary fixed point scheme** is the scheme stating that if $F: 2^{\mathbb{N}} \rightarrow 2^{\mathbb{N}}$ is arithmetical and

$$\forall X (X \subseteq F(X)),$$

then there is a w.o. (L, \prec) with max and sets $(X_\alpha : \alpha \in L)$ such that

$$\left\{ \begin{array}{l} X_{\min L} = \emptyset \\ X_{S(\alpha)} = F(X_\alpha); \\ X_\gamma = \bigcup_{\alpha \prec \gamma} X_\alpha \quad (\gamma \in \text{Lim}) \\ F(X_{\max L}) = X_{\max L}. \end{array} \right.$$

Stronger than $\Pi_1^1\text{-CA}_0$: does not require $X \subseteq Y \rightarrow F(X) \subseteq F(Y)$

Last slide!

Theorem (F-D S T Y)

The arithmetical inflationary fixed point scheme

\equiv *the maximal sequence principle*

\equiv *the closed orbit principle*

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Theorem (F-D S T Y)

The arithmetical inflationary fixed point scheme

\equiv *the maximal sequence principle*

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\mathcal{X}	f	V	Caristi	Ekeland
compact	continuous	continuous	WKL ₀	WKL ₀
	Baire 1	l.s.c.	ACA ₀	ACA ₀
arbitrary	continuous	continuous	ACA ₀	ACA ₀
	Baire	l.s.c.	TLPP ₀	Π_1^1 -CA ₀

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