

**Systems of propositions  
referring to each other:  
a model-theoretic view**

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## Theory $\equiv$

An important type of paradoxes combines negation and a kind of self-reference. The classical Liar paradox is a paradigmatic instance.

A similar paradox containing no (explicit) self-reference was proposed by Yablo about twenty years ago. This new paradox can be considered as an unfolding of the Liar paradox: it consists of propositions indexed by natural numbers such that each of the propositions states “all propositions with greater indices are wrong” .

*Remark.* The Yablo paradox generated plenty of papers (by Priest, Sorensen, Beall, Forster, and others) mainly discussed whether the paradox uses a kind of an implicit self-reference. This discussion is out of the subject of the talk.

*Our purpose:*

- (a) to investigate arbitrary systems of propositions some of them state that some of them (possibly, themselves) are wrong,
- (b) to establish which of these systems are paradoxical and which are not.

For this, we introduce a first-order theory  $\Xi$  of a language with one unary and one binary predicates,  $T$  and  $U$ , consisting of two easy axioms,  $A1$  and  $A2$ , which express natural relationships of propositions and their truth values.

**Definition.** Theory  $\Xi$ :

A1  $\forall xy (Tx \rightarrow (Uxy \rightarrow \neg Ty)),$

A2  $\forall x (\neg Tx \wedge \exists y Uxy \rightarrow \exists y (Uxy \wedge Ty)).$

Heuristically:

variables  $x, y$  mean propositions,

$Tx$  means “ $x$  is true”,

$Uxy$  means “ $x$  states that  $y$  is wrong”.

Thus:

A1 means “if  $x$  is true and states that  $y$  is wrong, then this  $y$  is indeed wrong”, while

A2 means “if  $x$  is wrong and states that some  $y$  are wrong, then some of these  $y$ 's witnesses that  $x$  is wrong”.

Note that the interpretation allows mute but wrong  $x$ 's.

What are properties of  $\Xi$ ?

**Theorem.**  $\Xi$  is a  $\Pi_2^0$  but not a  $\Sigma_2^0$  theory.

*Sketch of proof.*

1. Axiom  $A1$  is equivalent to the  $\Pi_1^0$  sentence

$$\forall xy (\neg Tx \vee \neg Uxy \vee \neg Ty)$$

and axiom  $A2$  is equivalent to the  $\Pi_2^0$  sentence

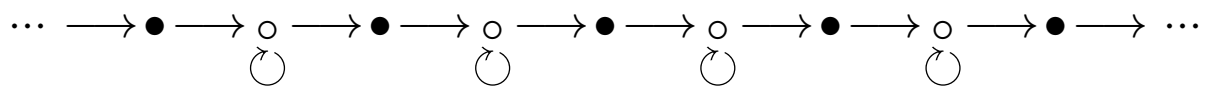
$$\forall xy \exists z (Tx \vee \neg Uxy \vee (Tz \wedge Uxz)).$$

Hence,  $\Pi_2^0$ .

Why not  $\Sigma_2^0$ ?

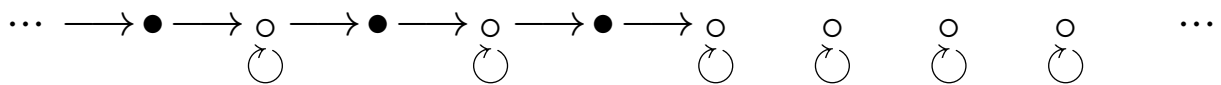
2. Let  $\mathfrak{A} = (X, T, U)$  where:  
 $X = \mathbb{Z}$  (the set of integers),  
 $T = \{2n + 1 : n \in \mathbb{Z}\}$ ,  
 $U = \{(n, n + 1), (2n, 2n) : n \in \mathbb{Z}\}$ .

We can picture it as follows:



(White nodes = elements  $x$  with  $\neg Tx$ ,  
black nodes = elements  $x$  with  $Tx$ ,  
arrows = pairs of elements  $x, y$  with  $Uxy$ .)

For any  $k \in \mathbb{Z}$ , let  $\mathfrak{A}_k$  be the submodel of  $\mathfrak{A}$  with the universe  $\{n \leq 2k : n \in \mathbb{Z}\} \cup \{2n : n \in \mathbb{Z}\}$ :



*Facts:*

1.  $\mathfrak{A}$  satisfies  $\Xi$ .
2.  $\mathfrak{A}_k$  do not satisfy  $\Xi$  and are isomorphic.

Toward a contradiction, assume that  $\Xi$  can be axiomatized by a set  $\Gamma$  of  $\Sigma_2^0$  sentences. So  $\mathfrak{A} \models \Gamma$ .

*Claim.* For any  $\gamma \in \Gamma$  there is  $k \in \mathbb{Z}$  with  $\mathfrak{A}_k \models \gamma$ .

[Why? Any sentence  $\gamma \in \Gamma$  has the form

$$\exists x_0 \dots x_m \forall y_0 \dots y_n \varphi(x_0, \dots, x_m, y_0, \dots, y_n)$$

with  $\varphi$  open. Since  $\mathfrak{A} \models \gamma$ , there are  $a_0, \dots, a_m \in X$  such that

$$\mathfrak{A} \models \forall y_0 \dots y_n \varphi(a_0, \dots, a_m, y_0, \dots, y_n).$$

Submodels preserve universal formulas, so any submodel of  $\mathfrak{A}$  satisfies the same whenever it has the  $a_0, \dots, a_m$ . Hence, if  $k = \max_{i \leq m} a_i$  then

$$\mathfrak{A}_k \models \varphi(a_0, \dots, a_m, y_0, \dots, y_n),$$

and so  $\mathfrak{A}_k \models \gamma$ .]

Now: all the  $\mathfrak{A}_k$  are isomorphic and so satisfy the same sentences. Therefore, each  $\mathfrak{A}_k$  should satisfy every  $\gamma \in \Gamma$ , and thus  $\Xi$ . A contradiction.

Hence, not  $\Sigma_2^0$ .



Other properties of  $\exists$ :

**Proposition.** *The theory  $\exists$ :  
has atomic saturated models of any cardinality,  
does not admit quantifier elimination,  
is neither complete nor model complete,  
not Horn,  
not categorical in any cardinality,  
undecidable.*

[E.g. why not categorical in cardinality 1?

There are single-point models, one with  $Tx$  and  
another with  $\neg Tx$ :

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The same for any cardinality.]

Let  $M(\Gamma)$  be the class of models of a theory  $\Gamma$ .

Model-theoretic operations which preserve or do not preserve  $\Xi$ :

**Proposition.**  $M(\Xi)$  is closed under:

*upper cones w.r.t.  $U$ ,*  
*disjoint sums and summands,*  
*unions of increasing chains,*  
*dualization,*  
*ultrafilter extensions,*  
*1-sandwichs.*

**Proposition.**  $M(\Xi)$  is not closed under:

*submodels,*  
*direct products and factors,*  
*bisimulations,*  
*images and pre-images of strong homomorphisms,*  
*non-empty intersections of two elementary sub-*  
*models,*  
*non-empty intersections of decreasing chains,*  
*2-sandwichs.*

## Collapses of models of $\Xi$

We describe a procedure that will be used for a natural classification of models of  $\Xi$ .

*Terminology.*

Let  $(X, T, U)$  and  $(X', T', U')$  be models of the language of  $\Xi$ . A map  $\pi : X \rightarrow X'$  is:  
a *quasi-isomorphism* w.r.t.  $T$  iff

$$\forall x (Tx \leftrightarrow T'\pi x),$$

a *bounded morphism* w.r.t.  $U$  iff

$$\forall xy (Uxy \rightarrow U'\pi x \pi y),$$

$$\forall xy' \exists y (U'\pi xy' \rightarrow \pi y = y' \wedge Uxy).$$

Note that bounded morphisms (aka p-morphisms) are stronger than strong homomorphisms.

**Lemma.** Let  $(X, T, U)$  and  $(X', T', U')$  be such that there exists  $\pi : X \rightarrow X'$  which is:  
 a surjection,  
 a quasi-isomorphism w.r.t.  $T$ ,  
 a bounded morphism w.r.t.  $U$ .  
 Then  $(X, T, U) \models \Xi$  iff  $(X', T', U') \models \Xi$ .

Informally:  $\pi$  “glues together” some of points with the same value of  $T$ .

After glueing *all* such points of a given model, it collapses to a model consisting of at most two points.

**Definition.** The *collapse* of a model  $(X, T, U)$  is the model  $(X', T', U')$  with the map  $\pi : X \rightarrow \{0, 1\}$  defined as follows:

$\pi x = 1$  if  $Tx$ , and  $\pi x = 0$  otherwise,

$X' = \{\pi x : x \in X\}$ ,

$T'x'$  iff  $x' = 1$  (i.e.  $T'1 \wedge \neg T'0$ ),

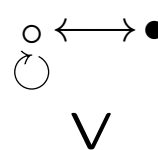
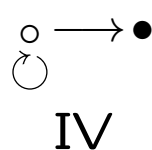
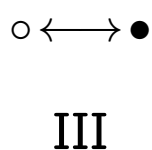
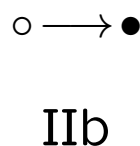
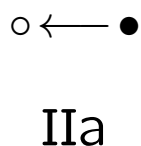
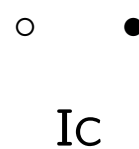
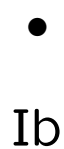
$U'x'y'$  iff  $\exists xy (\pi x = x' \wedge \pi y = y' \wedge Uxy)$ .

**Theorem.** *A model satisfies  $\Xi$  iff its collapse does. There exist exactly eight models that are the collapses of models of  $\Xi$ , listed below:*

	$X$	$T$	$U$
Ia	$\{0\}$	$\emptyset$	$\emptyset$
Ib	$\{1\}$	$\{1\}$	$\emptyset$
Ic	$\{0, 1\}$	$\{1\}$	$\emptyset$
IIa	$\{0, 1\}$	$\{1\}$	$\{(0, 1)\}$
IIb	$\{0, 1\}$	$\{1\}$	$\{(1, 0)\}$
III	$\{0, 1\}$	$\{1\}$	$\{(0, 1), (1, 0)\}$
IV	$\{0, 1\}$	$\{1\}$	$\{(0, 0), (0, 1)\}$
V	$\{0, 1\}$	$\{1\}$	$\{(0, 0), (0, 1), (1, 0)\}$

(Models with isomorphic  $U$  are denoted by the same Roman letters.)

The picture of the 8 models:



*Examples.*

1. Any model with empty  $U$  collapses to one of Ia, Ib, Ic.
2. The model  $\mathfrak{A}$  from the proof above (about complexity of  $\Xi$ ) collapses to V.

## Classification of models of $\Xi$

The class  $M(\Xi)$  is divided into eight subclasses consisting of the pre-images of eight models above under the collapse.

For  $X \in \{\text{Ia}, \text{Ib}, \text{Ic}, \text{IIa}, \text{IIb}, \text{III}, \text{IV}, \text{V}\}$ , let  $\Xi(X)$  denote the theory of pre-images of  $X$ .

Refining the theorem above:

### **Theorem.**

1.  $\Xi(\text{Ia}), \Xi(\text{Ib})$  are  $\Pi_1^0$  but not  $\Sigma_1^0$  theories.
2.  $\Xi(\text{Ic}), \Xi(\text{IIa}), \Xi(\text{IIb}), \Xi(\text{III}), \Xi(\text{IV})$  are  $\Delta_2^0$  but neither  $\Pi_1^0$  nor  $\Sigma_1^0$  theories.
3.  $\Xi(\text{V})$  is a  $\Pi_2^0$  but not  $\Sigma_2^0$  theory.

## Paradoxicality

**Definition.** A model  $(X, U)$  is *non-paradoxical* iff it can be expanded to some model  $(X, T, U)$  of  $\Xi$ , and *paradoxical* otherwise.

*Examples.*

1. A model of the Liar Paradox consists of one reflexive point:



and it is paradoxical.

2. A model of the Yablo Paradox is isomorphic to natural numbers with their usual ordering, and it is paradoxical.

3. A cycle is paradoxical iff its length is odd. (E.g. the Liar is a cycle of length 1.)

More examples can be obtained by considering specific classes of relations.



## Specific relations

**Proposition.** *If  $U$  is reflexive, then  $(X, U)$  is paradoxical.*

This generalizes the Liar Paradox, the model of which consists of one reflexive point.

Given  $(X, U)$ , a subset  $C$  of  $X$  is  *$U$ -cofinal* in  $X$  iff for each  $x \in X \setminus C$  there are  $y \in C$  and finitely many  $x_1, \dots, x_n \in X$  such that  $xUx_1U \dots Ux_nUy$ . E.g. if  $U^{-1}$  is well-founded, then the set of  $U$ -maximal elements is cofinal.

**Proposition.** *If  $U$  is transitive, then  $(X, U)$  is non-paradoxical iff the set of  $U$ -maximal elements is cofinal in  $X$ .*

This generalizes both the Liar and Yablo Paradoxes, the models of which are transitive without maximal elements.

**Proposition.** *If  $U^{-1}$  is well-founded, then  $(X, U)$  is non-paradoxical. Moreover, any  $(X, T, U)$  satisfying  $\Xi$  is uniquely determined by values of  $T$  on  $U$ -maximal elements.*

Consequently, if  $U^{-1}$  is well-founded and there is a unique  $U$ -maximal element  $x$  (e.g. if  $U^{-1}$  is extensional) then there are exactly two models  $(X, T, U)$  satisfying  $\Xi$ , depending on  $Tx$  or  $\neg Tx$ .

The following concept of hereditarily winning relations generalizes well-foundedness.

*A game of two players (Forster, Saveliev):*

Let  $(X, x_0, U)$  be a pointed model. The first player moves by choosing any  $x_1 \in X$  with  $Ux_1x_0$ ;

the second player (knowing about  $x_1$ ) moves by choosing any  $x_2 \in X$  with  $Ux_2x_1$ ;  
and so on.

A player *wins* iff he could choose an  $x_n \in X$  such that the play is stopped. (Possibly, nobody wins: e.g. look at a cycle.)

$(X, x_0, U)$  is *winning* for one of the players iff it admits a winning strategy for him.

$U$  is *hereditarily winning* iff  $(X \upharpoonright y, y, U \upharpoonright y)$  are winning for all  $y \in X$ .

*Fact:* Well-founded relations are hereditarily winning. There exist lots of non-well-founded hereditarily winning relations.

**Proposition.** *If  $U^{-1}$  is hereditarily winning, then  $(X, U)$  is non-paradoxical.*

**Main Theorem.** *The paradoxicality (and hence non-paradoxicality) is a  $\Delta_1^1$  but not elementary property.*

*Sketch of proof.*

1.  $(X, U)$  is non-paradoxical iff it satisfies the  $\Sigma_1^1$  formula

$$\exists T \left( T \text{ is a subset of } X \wedge \Xi^{X, T, U} \right).$$

2.  $(X, U)$  is non-paradoxical iff it satisfies the  $\Pi_1^1$  formula

$$\begin{aligned} \forall \pi \left( \pi \text{ is a map of } X \text{ into itself } \wedge |\text{ran } \pi| \leq 2 \right. \\ \rightarrow \left( \exists T' \subseteq \text{ran } \pi \right) \left( \exists U' \subseteq \text{ran } \pi \times \text{ran } \pi \right) \\ \Xi^{\text{ran } \pi, T', U'} \wedge \pi \text{ is a bounded morphism} \\ \left. \text{of } (X, U) \text{ onto } (\text{ran } \pi, U') \right). \end{aligned}$$

Here  $\pi$  collapses  $(X, U)$ , so  $(X, U)$  can be expanded to a model of  $\Xi$  by letting  $T = \pi^{-1}T'$ .

Hence,  $\Delta_1^1$ .

Why not elementary (= not first-order axiomatizable)?

3. Let  $x_{m,0} = x_{n,0}$  and  $x_{m,i} \neq x_{n,j}$  whenever  $(i, j) \neq (m, n)$ , and define:

$$X_n = \{x_{n,i}\}_{i \in \mathbb{Z}_{2n+1}}, \quad U_n = \{(x_{n,i}, x_{n,i+1})\}_{i \in \mathbb{Z}_{2n+1}},$$

$$X_\omega = \{x_{\omega,i}\}_{i \in \mathbb{Z}}, \quad U_\omega = \{(x_{\omega,i}, x_{\omega,i+1})\}_{i \in \mathbb{Z}},$$

$$X = \bigcup_{n < \omega} X_n, \quad U = \bigcup_{n < \omega} U_n,$$

$$X' = \bigcup_{n \leq \omega} X_n, \quad U' = \bigcup_{n \leq \omega} U_n.$$

Then:

- (i)  $(X, U)$  is paradoxical while  $(X', U')$  is not,
- (ii)  $(X, U) \prec (X', U')$  (an elementary submodel).

Hence, not elementary.

## Classification of non-paradoxical models

The class of non-paradoxical models is covered by five classes consisting of models that can be expanded to models of  $\Xi$  of classes I (a, b, or c, no matter), II (a or b), III, IV, and V. However the covering classes are *not disjoint*:

*Example.*

$* \longrightarrow *$

$* \longrightarrow *$

expansion:

$\circ \longrightarrow \bullet$

$\circ \longrightarrow \bullet$

$\circ \longrightarrow \bullet$

$\bullet \longrightarrow \circ$

collapse:

$\circ \longrightarrow \bullet$

$\circ \longleftrightarrow \bullet$

II

III

For  $X_1, \dots, X_n \in \{\text{I, II, III, IV, V}\}$ , let

$$X_1 + \dots + X_n$$

denote the class of non-paradoxical models such that for each  $i \in \{1, \dots, n\}$  the model can be expanded to a model of  $\Xi$  collapsed to  $X_i$ .

E.g. the model above is in  $\text{II} + \text{III}$ .

**Theorem.** *There exist exactly eleven disjoint classes of form  $X_1 + \dots + X_n$ , listed below:*

$$\text{I, II, III, IV, V,}$$

$$\text{II} + \text{III, II} + \text{IV, III} + \text{V, IV} + \text{V,}$$

$$\text{III} + \text{IV} + \text{V, II} + \text{III} + \text{IV} + \text{V.}$$