

# The Second Incompleteness Theorem for $\Sigma_1^0$ -semi-numerations

Albert Visser

Philosophy, Faculty of Humanities, Utrecht University

Wormshop



Moscow

October 20, 2017

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht

# Overview

## Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht

# Overview

Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



# Overview

Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



# Overview

Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



# Overview

## Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

### Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



# Numerations and Semi-numerations

We consider theories  $U$  in finite signature with arbitrarily complex axiom sets  $X$ . Let  $N : S_2^1 \triangleleft U$ . Consider any  $U$ -formula  $\alpha$ .

We say that  $\alpha$  *semi-numerates*  $X$  in  $N$  iff, for each  $A \in X$ , we have  $U \vdash \alpha(\underline{\Gamma A \neg})$ . Here we use  $N$ -numerals.

We say that  $\alpha$  *numerates*  $X$  in  $N$  iff, for each  $U$ -sentence  $A$ , we have  $A \in X$  iff  $U \vdash \alpha(\underline{\Gamma A \neg})$ .

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht

# Overview

Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

Preliminaries

**Example**

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht



# Example

Let  $Z$  be any set of numbers. By a theorem independently due to Feferman, Scott and Kripke, there is a  $\Sigma_1^0$ -predicate  $S^*(x)$  such that  $EA + \{S^*(\underline{n}) \mid n \in Z\} + \{\neg S^*(\underline{n}) \mid n \notin Z\}$  is consistent.

Consider the theory  $U$  axiomatized by  $X := EA + \{S^*(\underline{n}) \mid n \in Z\}$ . Let  $\beta$  be a standard representation of the finitely many axioms of EA. We take  $\zeta(x) := (\beta(x) \vee \exists y < x (x = \ulcorner S^*(\dot{y}) \urcorner \wedge S^*(y)))$ .

Clearly,

- ▶  $\zeta$  numerates  $X$  in  $U$ ,
- ▶  $\zeta$  is  $\Sigma_1^0$ ,
- ▶  $Z$  is 1-reducible to  $U$ .

So we have arbitrarily complex theories for which the axiom set is numerated by a  $\Sigma_1^0$ -formula (relative to the designated numbers).

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



# Overview

Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

Preliminaries

Example

**G2 for  $\Sigma_1^0$   
semi-numerations**

Example



# G2 for $\Sigma_1^0$ semi-numerations: stated

## Theorem

Suppose:

- ▶  $U$  is consistent,
- ▶  $U$  is axiomatized by  $X$ ,
- ▶  $N : EA \triangleleft U$ ,
- ▶  $\sigma$  is  $\Sigma_1^0$ ,
- ▶  $\sigma^N$  semi-numerates  $X$  in  $N$ .

Then,  $U \not\vdash (\diamond_{\sigma} \top)^N$ .

Alternative formulation of the conclusion: not  $N : U \triangleright (EA + \diamond_{\sigma} \top)$ .

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht

# G2 for $\Sigma_1^0$ semi-numerations: proved

## Proof.

Suppose  $X \vdash (\diamond_\sigma \top)^N$ . Let  $X_0$  be a finite subset of  $X$  such that  $X_0 \vdash EA^N$  and  $X_0 \vdash (\diamond_\sigma \top)^N$ . Let  $Y_0 := X_0 + \{\sigma^N(\ulcorner A \urcorner) \mid A \in X_0\}$ .

Using the fact that we have  $\Sigma_1^0$ -completeness in EA, we see that, for any  $B \in Y_0$ , we have  $Y_0 \vdash \Box_\sigma^N B$ .

Let  $\eta_0$  be a standard representation of the finite set  $Y_0$ . We find:  $Y_0 \vdash (\forall A \in \text{sent}_U (\Box_{\eta_0}^N A \rightarrow \Box_\sigma A))^N$ . **Note that we do not need internal  $\Sigma_1^0$ -collection since  $Y_0$  is standardly finite.**

It follows that  $Y_0 \vdash \diamond_{\eta_0}^N \top$ . But this contradicts the ordinary G2. ■

There is an alternative proof using Craig's trick. It seems that I need that alternative version to prove the corresponding variant of Feferman's Theorem.

Do we have the theorem also for  $S_2^1$ ?

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht

# Overview

Preliminaries

Example

G2 for  $\Sigma_1^0$  semi-numerations

Example

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



# Example: what the example does

The formalized  $\Delta_1^0$  axiom may fail in our setting. Thus the ordinary proof of G2 does not always work without an extra trick.

There is a  $\Sigma_1^0$ -formula  $\sigma(x)$  with the following properties.

- ▶  $\sigma$  defines the axioms of EA in the standard model, and, hence, numerates the axioms of EA in EA.
- ▶ EA knows that  $\sigma$  defines a finite set of axioms.
- ▶ EA knows that the theory defined by  $\sigma$  is between EA and  $EA + \Diamond_\beta \top$ , where  $\beta$  is a standard axiomatization of EA.
- ▶ EA does not prove formalized G2 for  $\sigma$ .

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



# Example: a quick peek at what it looks like 1

We work in EA. We define a Kripke model  $\mathcal{K}$ . Let  $p$  be any (possibly non-standard) number. Our model has nodes  $0, \dots, p+1$ . We set  $x \prec y$  iff  $x = 0$  and  $1 \leq y \leq p+1$ .

Let  $C$  be a single axiom for EA and let  $\beta(x) := x = \ulcorner C \urcorner$ .

We define the Solovay function  $h_p$  on  $\mathcal{K}$  for  $\beta$ .

- ▶  $l_p = 0$  iff  $\forall x h(x) = 0$ .
- ▶  $l_p = y$ , if  $0 \prec y$  and  $\exists x h(x) = y$ .
- ▶  $h_p(0) = 0$ ,
- ▶  $h_p(y+1) := \begin{cases} x & \text{if } h(y) \prec x \text{ and } \text{proof}_\beta(y, \ulcorner l_p \neq \dot{x} \urcorner) \\ h_p(y) & \text{otherwise} \end{cases}$

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



## Example: a quick peek at what it looks like 2

- ▶  $\text{def}(s, y)$  is the usual  $\Sigma_1^0$ -predicate for:  $s$  is a  $\Sigma_1^0$ -formula that defines  $x$ .
- ▶  $p$  is a partial term for the smallest inconsistency-proof of PA.
- ▶  $\sigma(x) :\leftrightarrow \beta(x) \vee (p \downarrow \wedge \exists y \leq p (x = \ulcorner \ell_p \neq (y + 1) \urcorner \wedge \exists s < p \text{ def}(s, y)))$ .
- ▶  $S^* :\leftrightarrow (p \downarrow \wedge \forall y \leq p \exists s < p \text{ def}(s, y))$ .

Note that  $\text{EA} + p \uparrow \vdash \forall x (\sigma(x) \leftrightarrow \beta(x))$  and  $\text{EA} \vdash \sigma(\ulcorner B \urcorner)$  iff  $\text{EA} \vdash \beta(\ulcorner B \urcorner)$ .

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht



# Example: a quick peek at what it looks like 3

## Selected Facts:

- ▶  $EA \vdash p \rightarrow \Box_{\beta} \neg S^*$ .
- ▶  $EA \vdash (y \leq x \wedge \Box_{\beta}(l_x \neq y + 1)) \rightarrow \Box_{\beta} \perp$ .
- ▶  $EA \vdash \bigvee_{y \leq x} (l_x = y + 1) \leftrightarrow \Box_{\beta} \perp$ .

## Theorem

- ▶  $EA \vdash \Box_{\sigma} \perp \leftrightarrow ((S^* \wedge \Box_{\beta} \Box_{\beta} \perp) \vee \Box_{\beta} \perp)$ .
- ▶  $EA \vdash \Box_{\sigma} \Diamond_{\sigma} \top \leftrightarrow (S^* \vee \Box_{\beta} \perp)$ .

Using methods due to Paris & Kirby, we can build a model  $\mathcal{N}$  of  $EA + S^* + \Diamond_{\beta} \Diamond_{\beta} \top$  and we are done.

Preliminaries

Example

G2 for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht

# Thank You



Preliminaries

Example

$G_2$  for  $\Sigma_1^0$   
semi-numerations

Example



Universiteit Utrecht